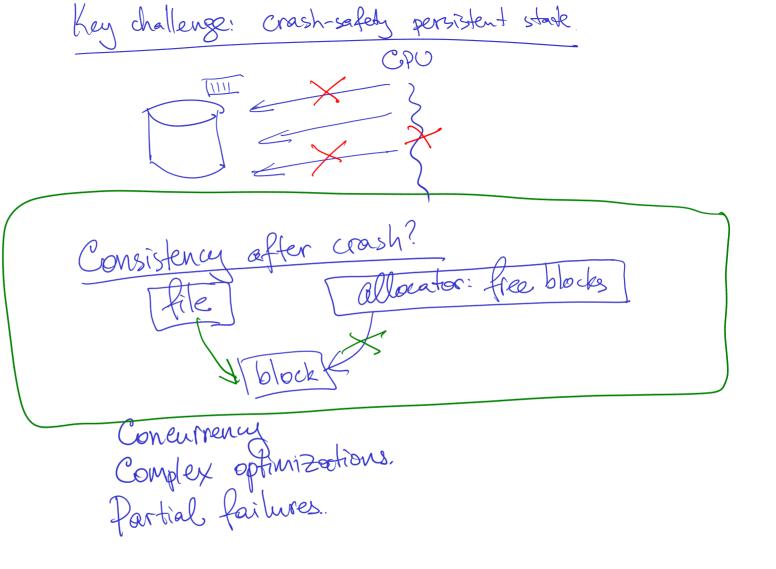
6.826: Systems + Verification. When is a system correct? Big ideas: locking, logging, consistency, security,... Principled view: Properties Specifications Proof

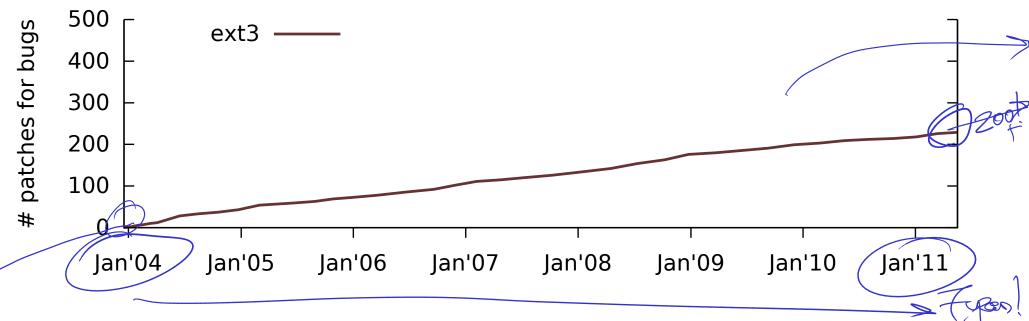
Motivation: begs. Impact Bugs Complex. memory safet Crash Concurrency deadlock lose deta Distributed Security =>races leak data aults Wrong result. Performance subtle Evolution logic

Example: fault tolerance Power failure -Disk died < Bad memory (Failure models Fail-stop "Byzantine" Contexts FS Zover DB Zover Distributed DB, KV SSD FTL



Example: file systems

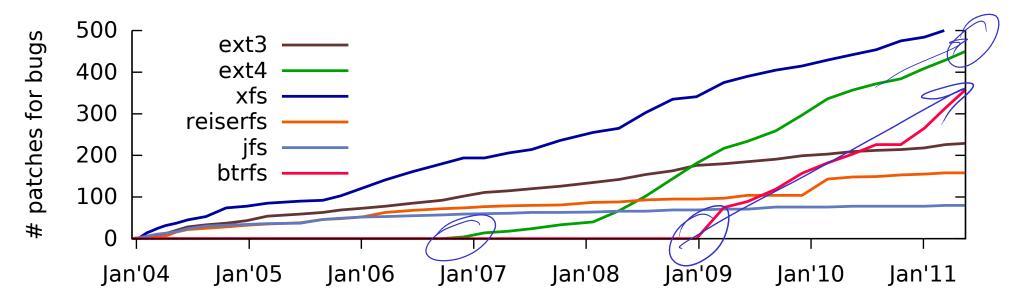
File systems are complex (e.g., Linux ext4 is \sim 60,000 lines of code) and have many bugs:



Cumulative number of patches for file-system bugs in Linux; data from [Lu et al., FAST'13]

Example: file systems

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New file systems (and bugs) are introduced over time

Write-ahead logging Order BARRIER Log: (addr, block) Data -> BARRIER Har batched. -> BARRIER Optimization 3: disk buffering. BARRIER. Optimization 2: log-sypass writes. READ (addr) sblock Optimitation 3: checksum logging. [Hdr=chedcsum (log). BARRIER (all) Jok. CRC 32 SHA-256. (P)

Bug example: data corruption

```
commit 353667d8ced4dc53281c88150ad295e24bc4b4c5
--- a/f/jbd//heckpoint.c
+++ b/fs/ibd/checkpoint.c
@@ -504,7 +503,25 @@ int cleanup_journal_tail(journal_t *journal)
             spin_unlock(&journal->j_state_lock);
             return 1;
     }
     spin_unlock(&journal->j_state_lock);
+
+
     /*
+
      * We need to make sure that any blocks that were recently written out
      * --- perhaps by log_do_checkpoint() --- are flushed_out before we
      * drop the transactions from the journal. It's unlikely this will be
      * necessary, especially with an appropriately sized journal, but we
+
       need this to guarantee correctness. Fortunately
+
       cleanup_journal_tail() doesn't get called all that often.
+
      */
+
     if (journal->j_flags & JFS_BARRIER)
+
             blkdev_issue_flush(journal->j_fs_dev, GFP_KERNEL, NULL);
+
     spin_lock(&journal->j_state_lock);
+
    if (!tid_gt(first_tid, journal->j_tail_sequence)) {
+
             spin_unlock(&journal->j_state_lock);
+
             /* Someone else cleaned up journal so return 0 */
+
             return 0:
+
+
     }
     /* OK, update the superblock to recover the freed space.
      * Physical blocks come first: have we wrapped beyond the end of
      * the log? */
```

Bug example: data disclosure

- Two optimizations in Linux ext4: direct data write and log checksum
- Subtle interaction: new file can contain other users' data after crash
- Bug introduced in 2008, fixed in 2014 (six years later!)

Author: Jan Kara <jack@suse.cz> Date: Tue Nov 25 20:19:17 2014 -0500

ext4: forbid journal_async_commit in data=ordered mode

Option journal_async_commit breaks gurantees of data=ordered mode as it sends only a single cache flush after writing a transaction commit block. Thus even though the transaction including the commit block is fully stored on persistent storage, file data may still linger in drives caches and will be lost on power failure. Since all checksums match on journal recovery, we replay the transaction thus possibly exposing stale user data.

[...]

Avoiding bugs? Testing. Effective. Cannot ensure & bugs. Model checking. Exploring all possible states. State explosion Verification Prove Ø bugs. Active area. Tools: Cog, Lean, Z3, CVC

Ideas: Refinement, Sep. logic,... Demand: Security, complex. Successes: Amazon, Complert, Chrome / Firefox coupts.

Variations. No spec. -> lightweight implied spec. No proof -> verifier automates proof. Vo exec. code No verifier.

ogistices Papers. > Cog. > Collab decussion. > Individual solution. abs. = Grades: Labs Summaries. Class part. Butler Lampson. Nöckolai Tei Steller